Memory Segmentation and Physical address calculation

Introduction

Memory segmentation is nothing which is the methods where whole memory is divided into the smaller parts. in 8086 microprocessor memory are divided into for parts which is known as the segments. these segments are data segment, code segment, stack segment and extra segment.

Memory Segmentation

The total memory size is divided into segments of various sizes.

A segment is just an area in memory.
The process of dividing memory this way is called Segmentation.

Memory Segmentation

- In memory, data is stored as bytes.
 - Each byte has a specific address.
 - Intel 8086 has 20 lines address bus.
 - With 20 address lines, the memory that can be addressed is 2 power20 bytes.
 - 2power20= 1,048,576 bytes (1 MB).
 - 8086 can access memory with address

ranging from 00000 H to FFFFF H.

segments.

In 8086, memory has four different types of segments.

These are:

- Code Segment
- Data Segment
- Stack Segment
- Extra Segment

Segment Register

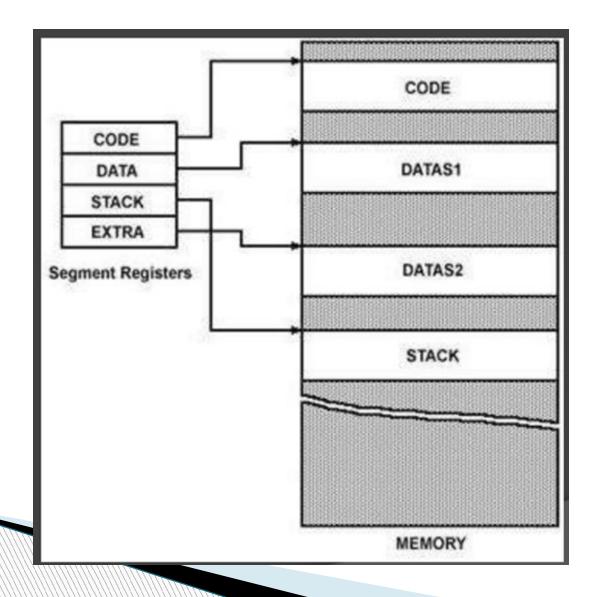
bits

⁷⁷ Each of these segments are addressed by an address stored in corresponding segment register.

These registers are 16-bit in size. Each register stores the base address (starting address) of the corresponding segment.

⁷⁷ Because the segment registers cannot store 20 bits, they only store the upper 16

Segment Register



How is a 20-bit address obtained if there are only 16bit registers?

The answer lies in the next few slides.

⁷⁷ The 20-bit address of a byte is called its Physical Address.

¹⁷ But, it is specified as a Logical Address.

Togical address is in the form of:

Base Address : Offset

⁷⁷ Offset is the displacement of the memory location from the starting location of the segment.

Example

The value of Data Segment Register (DS) is 2222 H.

To convert this 16-bit address into 20-bit, the BIU appends 0H to the LSBs of the address.

⁷⁷ After appending, the starting address of the Data Segment becomes 22220H.

If the data at any location has a logical address specified as:

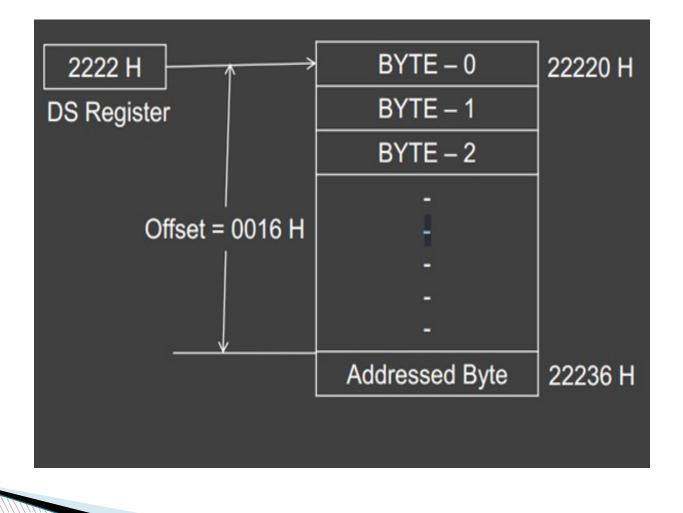
2222 H : 0016 H

Then, the number 0016 H is the offset. $_{77}$ 2222 H is the value of DS.

- To calculate the effective address of the memory, BIU uses the following formula: ⁷¹ Effective Address = Starting Address of Segment + Offset
 - To find the starting address of the segment, BIU appends the contents of Segment Register with 0H.
 - Then, it adds offset to it.

Therefore: EA = 22220 H+ 0016 H

22236 H



Maximum size of segment

- All offsets are limited to 16-bits.
- It means that the maximum size
- possible for segment is 2power16= 65,535
- bytes (64 KB).

- The offset of the first location within the
- segment is 0000 H.
- The offset of the last location in the segment is FFFF H

Where to look offset

Segment	Offset Registers	Function
CS	IP	Address of the next instruction
DS	BX, DI, S	Address of data
SS	SP, BP	Address in the stack
ES	BX, DI, SI	Address of destination data (for string operations)

Example

- The contents of the following registers are:
- ► CS = 1111 H
- ► DS = 3333 H
- ► SS = 2526 H
- ► IP = 1232 H
- ► SP = 1100 H
- ► DI = 0020 H
- Calculate the corresponding physical addresses for the address bytes in CS, DS and SS.

Solution

1. CS = 1111 H

The base address of the code segment is 11110 H.

- Effective address of memory is given by 11110H + 1232H = 12342H.
- 2. DS = 3333 H

The base address of the data segment is 33330 H.

- Effective address of memory is given by 33330H + 0020H = 33350H.
- 3. SS = 2526 H

The base address of the stack segment is 25260 H.

Effective address of memory is given by 25260H + 1100H =

263501

Scope of Research

Here we can various methods of memory segmentation, or we can say that logical segmentations. In logical segmentation we can divide the memory in logically. And in physical segmentation divide the memory physically.